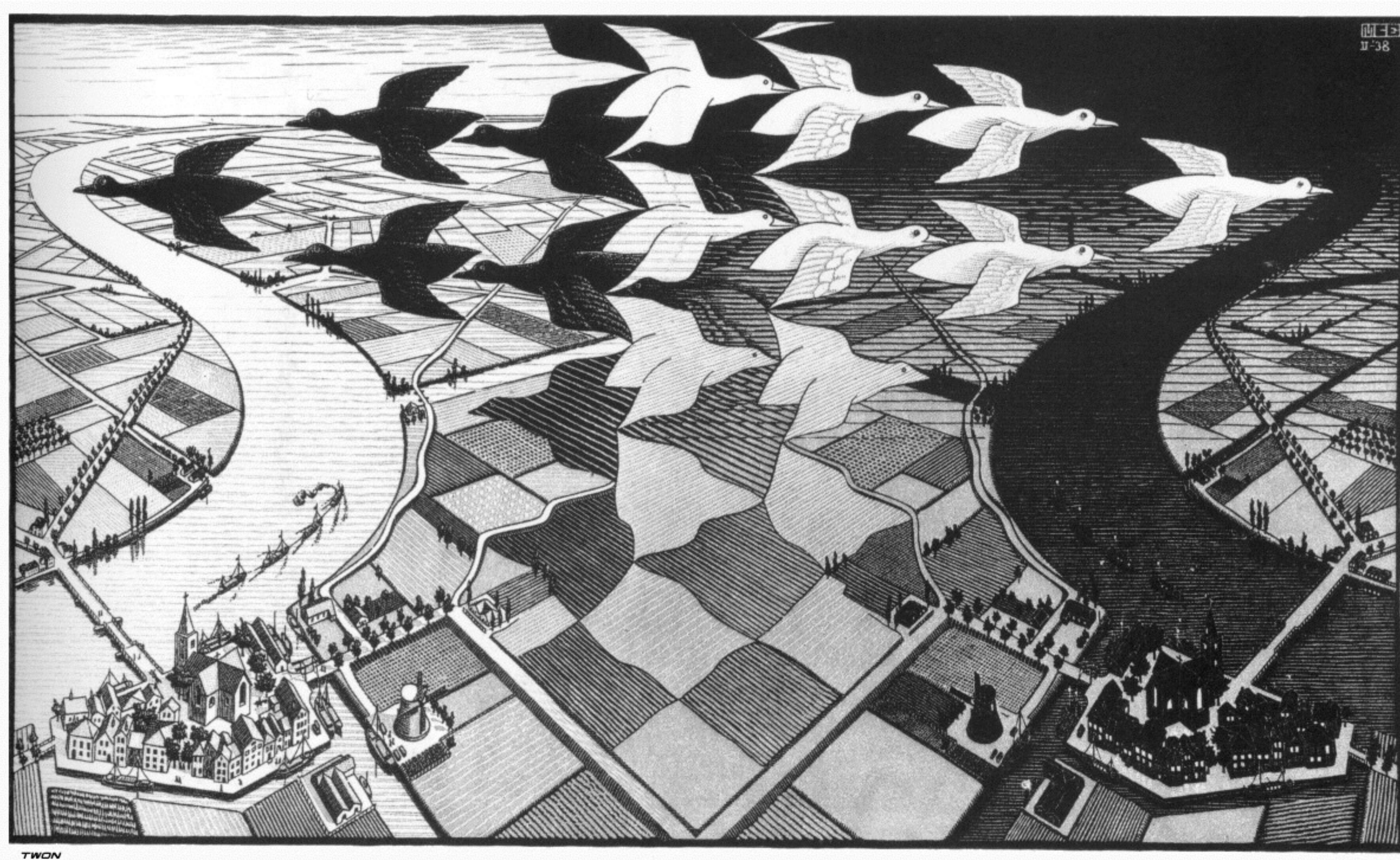


The Illusion Of Execution





Thanks!



Nitsan Wakart (@nitsanw)
Lead Performance Engineer, Azul Systems



PSYCHOSOMATIC, LOBOTOMY, SAW

It's X, you'll need Y, I'll get Z

Monday, 1 December 2014

The Escape of ArrayList.iterator()

{This post assumes some familiarity with JMH. For more JMH related content start at the new and improved **JMH Resources Page** and branch out from there!}

Escape Analysis was a much celebrated optimisation added to the JVM in Java 6u23:

"Based on escape analysis, an object's escape state might be one of the following:

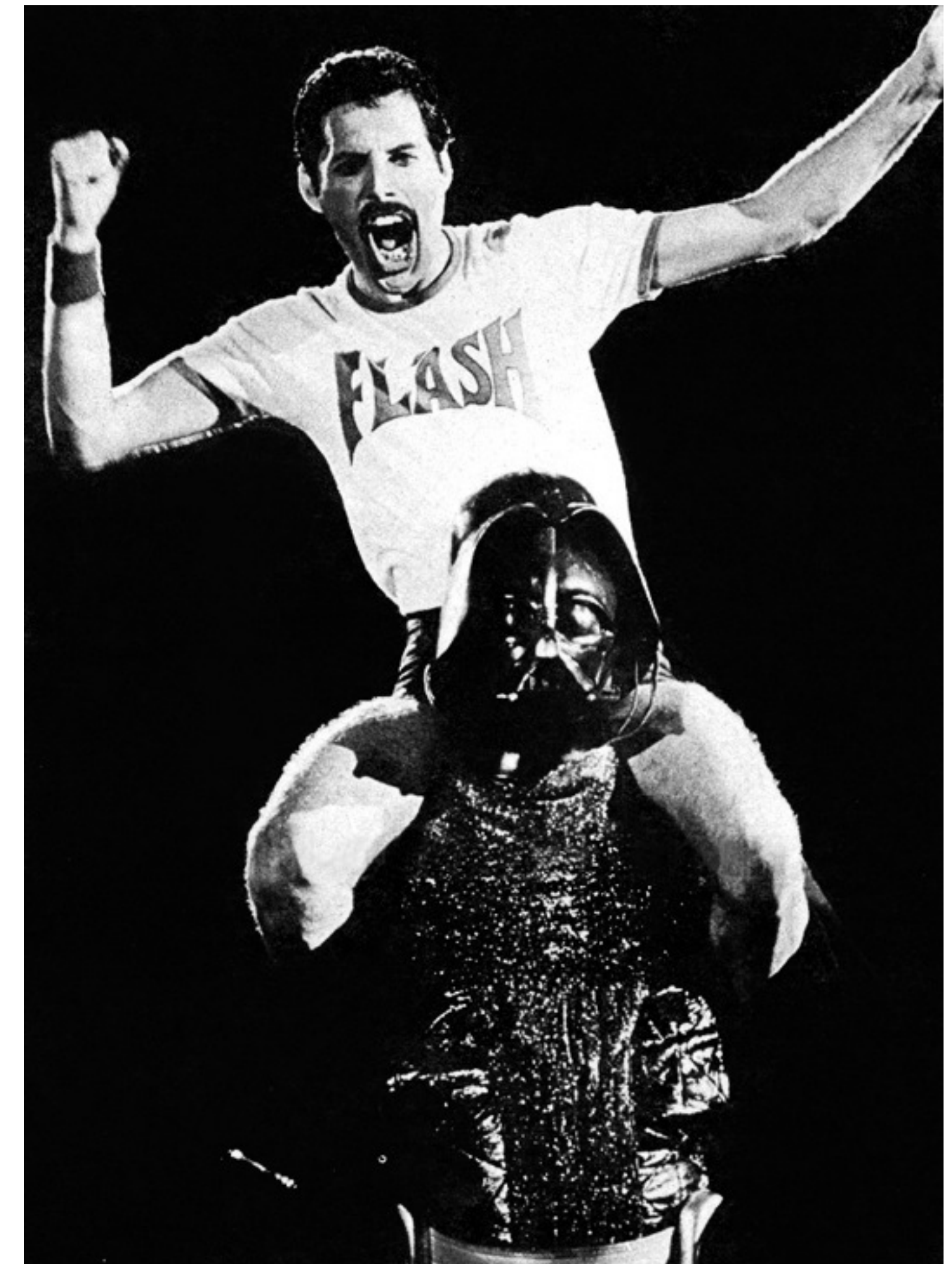
- **GlobalEscape** – An object escapes the method and thread. For example, an object stored in a static field, or, stored in a field of an escaped object, or, returned as the result of the current method.
- **ArgEscape** – An object passed as an argument or referenced



<http://psy-lob-saw.blogspot.com>

The JVM is a Magical Place

- You write some Java/Scala/Clojure/....
- Compile/Pack/Deploy
- ... MAGIC! ...
- JVM executes the “code”



The Way Down

Your (byte)code...

Is executed (interpreter/C1/C2)...

by a JVM (OpenJDK/Oracle/IBM/Zing)...

Which is a process of an OS (specific native libraries)...

Which is running on some hardware (specific instruction set)...

“I fucking hate assembly”

–Daniel San



All the assembly you need (for this talk)

- Each line is a single instruction
- Instructions take 1 to N CPU cycles
- MOV DST, SRC (this is Intel syntax)
- CMP A,B + JE DST -> if(A == B) goto DST
- INC A -> A++;
- **DON'T PANIC!!!!**

Code -> Assembly

```

public class Holder {
    T hold;
    int setCount;
    boolean set(T love) {
        if(love == hold) {
            return false;
        }
        else if (love != null &&
            love.equals(hold)) {
            return false;
        }
        hold = love;
        setCount++;
        return true;
    }
}

```

```

0x0000000010b065e8a: inc     DWORD PTR [rsi+0xc] ; *putfield setCount 0x0000000010b065f11: mov     rbp,rsi
; - Holder::set@39 (line 15)
0x0000000010b065e8d: mov     r10,rsi
0x0000000010b065e90: mov     r11,rdx
0x0000000010b065e93: shr     r11,0x3
0x0000000010b065e97: mov     DWORD PTR [rsi+0x10],r11d
; - Holder::set@39 (line 15)
0x0000000010b065e9b: shr     r10,0x9
0x0000000010b065e9f: mov     BYTE PTR [r11+r10*1],r10b
; - Holder::set@39 (line 15)
0x0000000010b065e99: mov     BYTE PTR [r11+r10*1],r10b
; - Holder::set@39 (line 15)
0x0000000010b065e20: mov     r10d,DWORD PTR [rsi+0x8]
; *putfield hold
0x0000000010b065e24: shl     r10,0x3
; - Holder::set@29 (line 14)
0x0000000010b065e28: cmp     rax,r10
; *getfield hold
0x0000000010b065e2b: jne     0x0000000010b065ead ; {runtime_call}
; - Holder::set@2 (line 8)
0x0000000010b065e31: data32 0x0000000010b065eb2: add     rsp,0x30
0x0000000010b065e34: nop
; *invokevirtual equals
0x0000000010b065e3c: data32 0x0000000010b065eb6: pop     rbp
; - Holder::set@19 (line 11)
0x0000000010b065e37: test   DWORD PTR [rip+0xffffffff71c143],eax
; {poll_return}
[Verified Entry Point]
0x0000000010b065e40: mov     DWORD PTR [rsp-0x14000],eax
0x0000000010b065e47: push   rbp
; *synchronization entry
0x0000000010b065e48: sub     rsp,0x30
; - Holder::set@-1 (line 8)
0x0000000010b065e4c: mov     r10d,DWORD PTR [rsi+0x10]
; - Holder::set@19 (line 11)
0x0000000010b065e4c: mov     esi,0xffffffffde
; *getfield hold
0x0000000010b065e50: mov     r11,r10
; - Holder::set@2 (line 8)
0x0000000010b065e53: shl     r11,0x3
; *invokevirtual equals
0x0000000010b065e57: xor     eax,eax
; - Holder::set@19 (line 11)
0x0000000010b065e59: cmp     rdx,r11
; {runtime_call}
0x0000000010b065e5d: jle     0x0000000010b065eb2 ; *if_acmpne
; - Holder::set@5 (line 8)
0x0000000010b065e5e: mov     r8d,DWORD PTR [rsi+0x8]
; implicit exception: dispatches to
; - java.lang.Integer::equals@1 (line 764)
0x0000000010b065f11: mov     rbp,rsi
; - Holder::set@19 (line 11)
0x0000000010b065f12: cmp     r8d,0xef5d495b
; {oop('java/lang/Integer')}
; - Holder::set@19 (line 11)
0x0000000010b065f19: jne     0x0000000010b065ee0 ; *invokevirtual equals
; - Holder::set@19 (line 11)
0x0000000010b065f19: mov     esi,0xffffffffde
; - Holder::set@19 (line 11)
0x0000000010b065f19: call   0x0000000010b038f20 ; OopMap{rbp=Oop [8]=Oop [20]=NarrowOop off=212}
; implicit exception: dispatches to
; *instanceof [Stub Code]
; - java.lang.Integer::equals@1 (line 764)
0x0000000010b065f40: jmp     0x0000000010b05f0a0 ; {no_reloc}
; - Holder::set@19 (line 11)
; {runtime_call}
0x0000000010b065f45: call   0x0000000010b065f4a
; {runtime_call}
0x0000000010b065f4a: sub     QWORD PTR [rsp],0x5
0x0000000010b065f4f: jmp     0x0000000010b038b00 ; {runtime_call}
; - Holder::set@19 (line 11)
0x0000000010b065efc: mov     QWORD PTR [rsp+0x8],rdx
; - Holder::set@19 (line 11)
0x0000000010b065f01: mov     esi,0xfffffffff4
; *instanceof
; - java.lang.Integer::equals@12 (line 765)
0x0000000010b065f06: nop
; - Holder::set@19 (line 11)
0x0000000010b065f07: call   0x0000000010b038f20 ; OopMap{rbp=Oop [8]=Oop [20]=NarrowOop off=236}
; - Holder::set@19 (line 11)
; *instanceof
; - java.lang.Integer::equals@1 (line 764)
0x0000000010b065e81: mov     r11d,DWORD PTR [r10+0xc]
; - Holder::set@19 (line 11)
0x0000000010b065e85: cmp     r8d,r11d
; - Holder::set@19 (line 11)
0x0000000010b065e88: je      0x0000000010b065eb2 ; *if_icmpne
; - Holder::set@19 (line 11)
; {runtime_call}
0x0000000010b065f0e: call   0x0000000010b038f20 ; OopMap{rbp=Oop [8]=Oop [16]=NarrowOop off=134}
; - Holder::set@19 (line 11)
; *instanceof #1
; - java.lang.Integer::equals@1 (line 764)
; - Holder::set@19 (line 11)
; {runtime_call}
OopMap{rbp=Oop [8]=Oop [20]=NarrowOop off=212}
; - Holder::set@19 (line 11)
; {runtime_call}
OopMap{rbp=Oop [8]=Oop off=236}

```

How to print assembly: psy-lob-saw.blogspot.com/2013/01/java-print-assembly.html

‘Focus!’

Master Miagi



Method Prologue

```
# {method} 'set' '(Ljava/lang/Object;)Z' in 'Holder'
# this:      rsi:rsi      = 'Holder'
# parm0:     rdx:rdx     = 'java/lang/Object'
#          [sp+0x40] (sp of caller)
0x0000000010bec3aa0: mov     r10d,DWORD PTR [rsi+0x8]
0x0000000010bec3aa4: shl     r10,0x3
0x0000000010bec3aa8: cmp     rax,r10
0x0000000010bec3aab: jne     0x0000000010be97960 ; {runtime_call}
0x0000000010bec3ab1: data32 xchg ax,ax
0x0000000010bec3ab4: nop     DWORD PTR [rax+rax*1+0x0]
0x0000000010bec3abc: data32 data32 xchg ax,ax
[Verified Entry Point]
0x0000000010bec3ac0: mov     DWORD PTR [rsp-0x14000],eax
0x0000000010bec3ac7: push   rbp
0x0000000010bec3ac8: sub     rsp,0x30 ;*synchronization entry
; - Holder::set@-1 (line 6)
```

Assuming hold != love

```
mov     r10d,DWORD PTR [rsi+0x10] ;*getfield hold/Holder::set@2 (line 8)
mov     r11,r10
shl     r11,0x3
xor     eax,eax
cmp     rdx,r11
je      0x000000010b065eb2 ; Holder::set@5 (line 8)
mov     r8d,DWORD PTR [rdx+0x8] ; implicit exception: dispatches to 0x000000010b065f11
cmp     r8d,0xef5d495b ; {oop('java/lang/Integer')}
jne     0x000000010b065ebe ;*invokevirtual equals/Holder::set@(line 11)
mov     r8d,DWORD PTR [r12+r10*8+0x8] ; implicit exception: dispatches to 0x000000010b065ef9
cmp     r8d,0xef5d495b ; {oop('java/lang/Integer')}
jne     0x000000010b065edd ;*instanceof/Integer::equals@(line 764)/Holder::set@(line 11)
mov     r8d,DWORD PTR [rdx+0xc]
shl     r10,0x3 ;*checkcast/Integer::equals@(line 765)/Holder::set@(line 11)
mov     r11d,DWORD PTR [r10+0xc]
cmp     r8d,r11d
je      0x000000010b065eb2 ; - java.lang.Integer::equals@18 (line 765)/Holder::set@(line 11)
inc     DWORD PTR [rsi+0xc] ;*putfield setCount/Holder::set@(line 15)
mov     r10,rsi
mov     r11,rdx
shr     r11,0x3
mov     DWORD PTR [rsi+0x10],r11d
shr     r10,0x9
movabs  r11,0x10a4e9000
mov     BYTE PTR [r11+r10*1],r12b ;*putfield hold/Holder::set@(line 14)
mov     eax,0x1 ;*getfield hold/Holder::set@(line 8)
add     rsp,0x30; This is 0x000000010b065eb2
pop     rbp
test    DWORD PTR [rip+0xfffffffffff71c143],eax # 0x000000010a782000; {poll_return}
ret
```

Conditional Inlining

```
mov    r8d,DWORD PTR [rdx+0x8] ; implicit exception: dispatches to 0x000000010b065f11
cmp    r8d,0xef5d495b          ; {oop('java/lang/Integer')}
jne    0x000000010b065ebe      ;*invokevirtual equals/Holder::set@(line 11)
mov    r8d,DWORD PTR [r12+r10*8+0x8] ; implicit exception: dispatches to 0x000000010b065ef9
cmp    r8d,0xef5d495b          ; {oop('java/lang/Integer')}
jne    0x000000010b065edd      ;*instanceof/Integer::equals@(line 764)/Holder::set@(line 11)
mov    r8d,DWORD PTR [rdx+0xc]
shl    r10,0x3                 ;*checkcast/Integer::equals@(line 765)/Holder::set@(line 11)
mov    r11d,DWORD PTR [r10+0xc]
cmp    r8d,r11d
je     0x000000010b065eb2      ; - Integer::equals@(line 765)/Holder::set@(line 11)
```

- The JIT compiler inlined Integer::equals
- Optimistically compiled, but checks class to trigger deoptimization
- Our code is compiled based on it's usage!

Compiler Reordering

```
inc    DWORD PTR [rsi+0xc] ;*putfield setCount/Holder::set@(line 15)
mov    r10,rsi
mov    r11,rdx
shr    r11,0x3
mov    DWORD PTR [rsi+0x10],r11d
shr    r10,0x9
movabs r11,0x10a4e9000
mov    BYTE PTR [r11+r10*1],r12b ;*putfield hold/Holder::set@(line 14)
```

- The JIT compiler reordered lines 15 and 14
- This can happen to YOU!

Instructions you have?
Get cycles, you must!
Hmmm, yes.....



```
$jstack 12345
```

How many threads are running?

Multi-Tasking OS

- More processes than cores
 - Yay!!! Gimme ALL the threads!
 - ... unless they all want to run at the same time
- Scheduling and interrupts
 - Fairness (thread priority, starvation)
 - Context Switching (overhead)

Thread States

Waiting - Thread is blocked (IO/wait()/park()...)

Ready - You are queued for execution (yield)

Running - CPU is busy with your logic

Have you ever written a single threaded Java application?

The JVM Process: Threads

- How many threads for an application?
 - Application Java threads (Main, Thread etc)
 - Application native threads (native lib)
 - JVM Threads (GC, Compiler, JMX, RMI...)
- There's no such thing as a single threaded Java application

Threads Example

(Oracle JDK8u20, on i5/dual core laptop, no args)

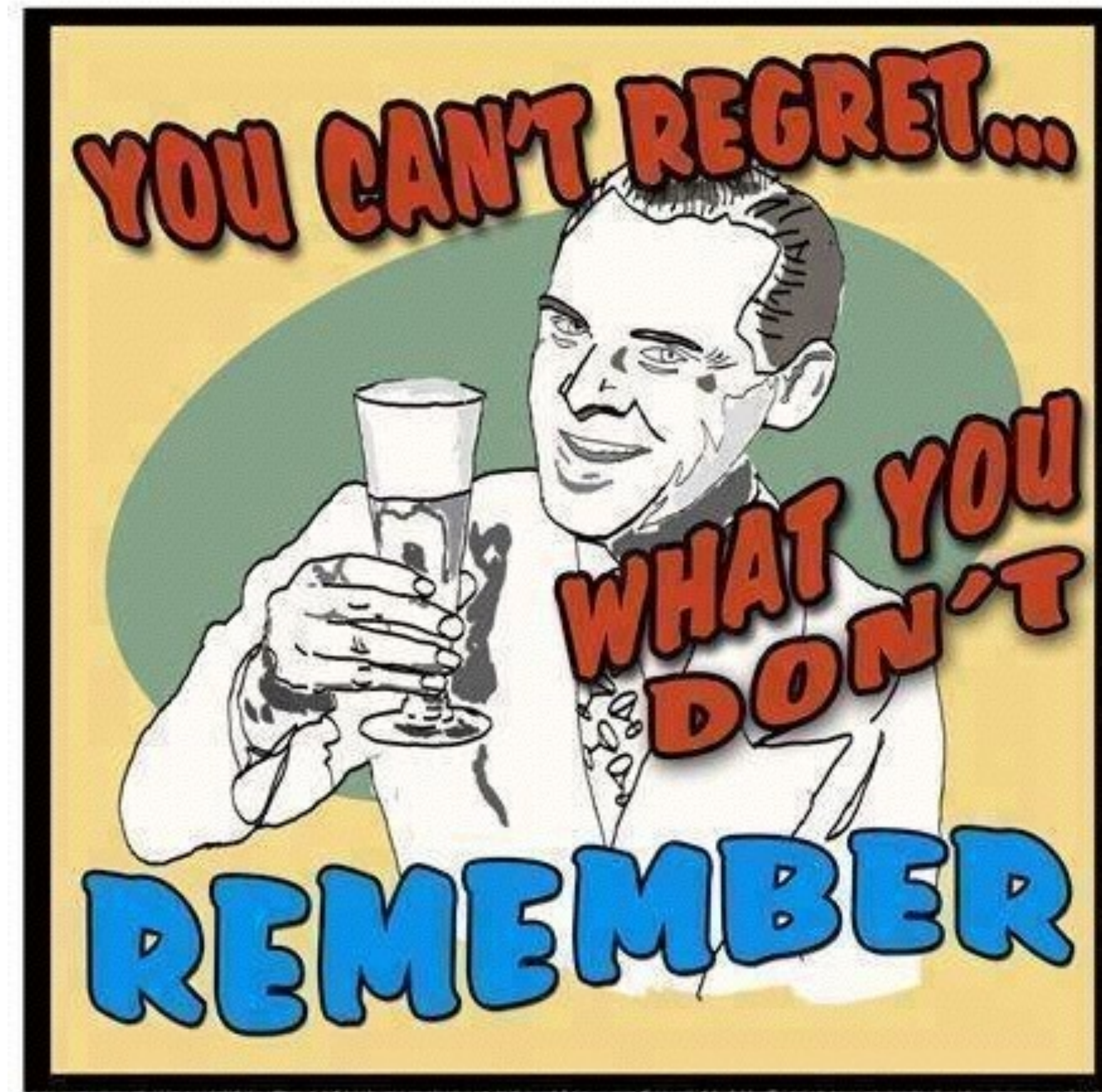
- Application threads
- 4 GC Threads (ParallelGC)
- 3 Compiler threads (1 C1 + 2 C2)
- ... and others
- 15 threads reported by jstack, 19 by OS

So What?

- Try and avoid having more **ready/running** threads than cores
 - Look out for long run queue (avg. load), ctx switches
- Consider controlling OS resource management
 - Use taskset/numactl/isocpus to reduce contention
 - Note that JVM is not aware of taskset/numactl/isocpus
- Consider controlling JVM thread counts



The Memory Illusion



How long to
`pub.getBeer()`?

Why is MOV important?

```
mov r10d,DWORD PTR [rsi+0x10] ;*getfield hold/Holder::set@2 (line 8)
```

```
mov r11,r10
```

```
shl r11,0x3
```

```
xor eax,eax
```

```
cmp rdx,r11
```

```
je 0x000000010b065eb2 ; Holder::set@5 (line 8)
```

```
mov r8d,DWORD PTR [rdx+0x8] ; implicit exception: dispatches to 0x000000010b065f11
```

```
cmp r8d,0xef5d495b ; {oop('java/lang/Integer')}
```

```
jne 0x000000010b065ebe ;*invokevirtual equals/Holder::set@(line 11)
```

```
mov r8d,DWORD PTR [r12+r10*8+0x8] ; implicit exception: dispatches to 0x000000010b065ef9
```

```
cmp r8d,0xef5d495b ; {oop('java/lang/Integer')}
```

```
jne 0x000000010b065edd ;*instanceof/Integer::equals@(line 764)/Holder::set@(line 11)
```

```
mov r8d,DWORD PTR [rdx+0xc]
```

```
shl r10,0x3 ;*checkcast/Integer::equals@(line 765)/Holder::set@(line 11)
```

```
mov r11d,DWORD PTR [r10+0xc]
```

```
cmp r8d,r11d
```

```
je 0x000000010b065eb2 ; - java.lang.Integer::equals@18 (line 765)/Holder::set@(line 11)
```

```
inc DWORD PTR [rsi+0xc] ;*putfield setCount/Holder::set@(line 15)
```

```
mov r10,rsi
```

```
mov r11,rdx
```

```
shr r11,0x3
```

```
mov DWORD PTR [rsi+0x10],r11d
```

```
shr r10,0x9
```

```
movabs r11,0x10a4e9000
```

```
mov BYTE PTR [r11+r10*1],r12b ;*putfield hold/Holder::set@(line 14)
```

```
mov eax,0x1
```

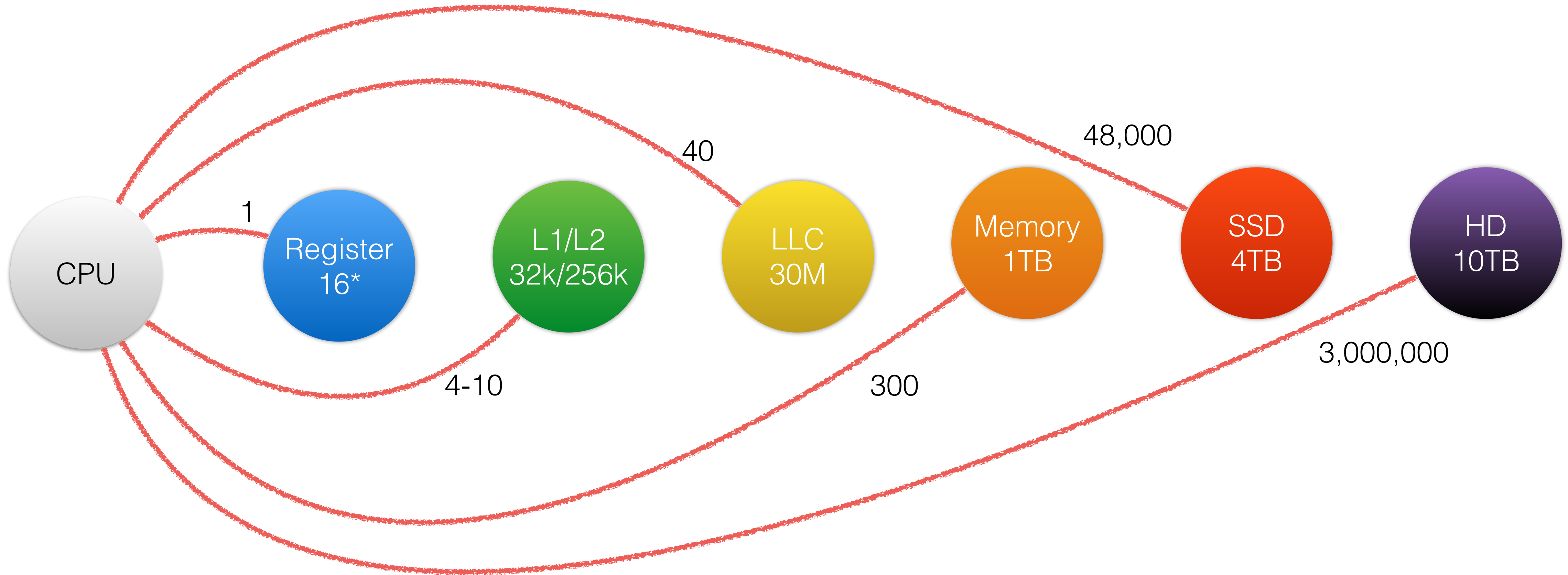
```
add rsp,0x30; This is 0x000000010b065eb2
```

```
pop rbp
```

```
test DWORD PTR [rip+0xffffffff71c143],eax # 0x000000010a782000; {poll_return}
```

```
ret
```

Memory Topology



Converting to cycles, assumed 3 cycles per nano-second
http://www.eecs.berkeley.edu/~rcs/research/interactive_latency.html

Beer Cache Hierarchy

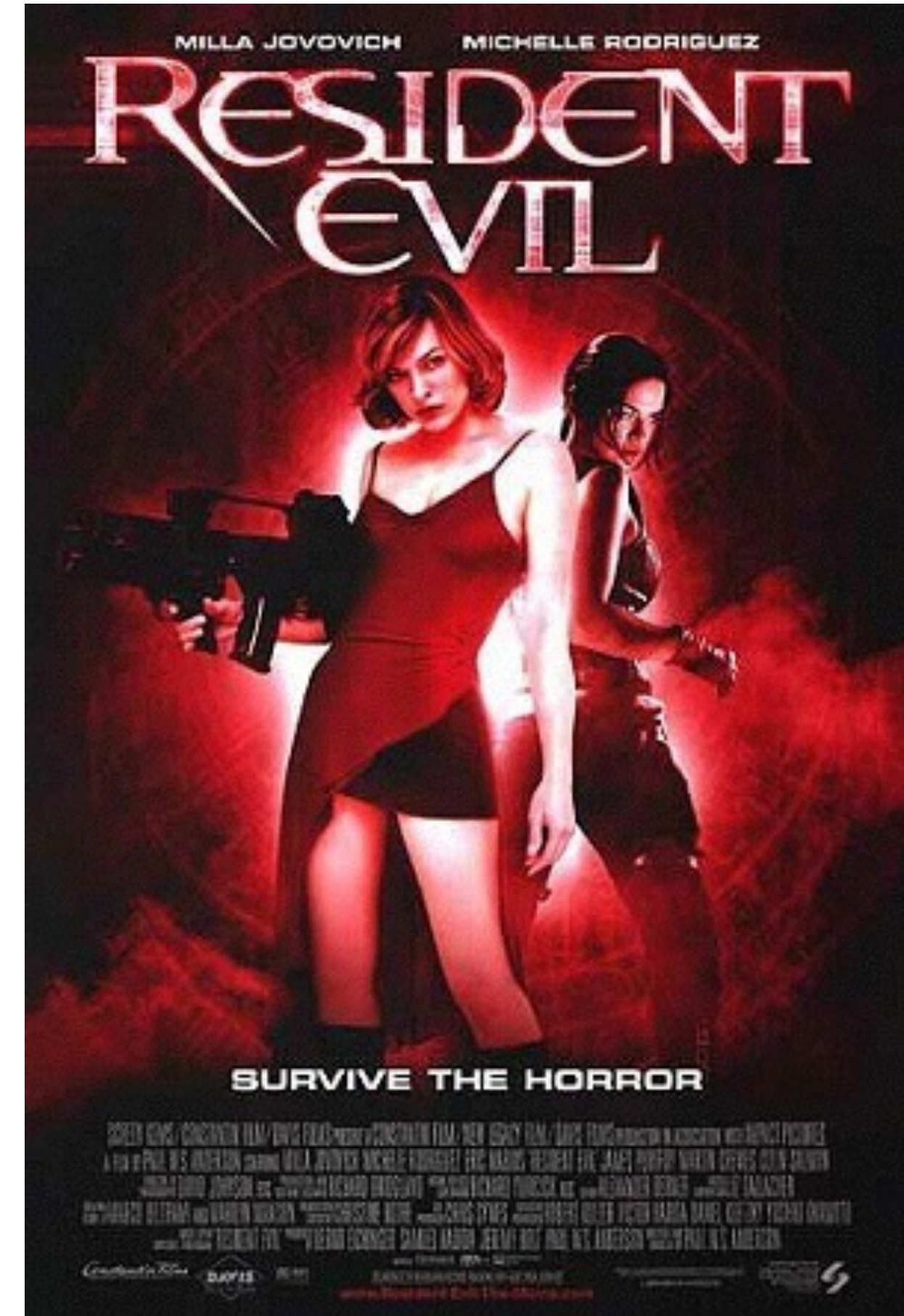
- Register - Bottle at your lips (1 sec)
- L1/L2 - ... on table (4 - 10 secs)
- LLC - ... waiting on the bar (40 secs)
- Main memory - ... at store across the road (5 mins)
- SSD - ... in another country? (13.3 hrs)
- HD - ... on the moon? (34 days, 17 hrs)

See: <http://blog.netopyr.com/2014/11/28/reactions-to-the-beer-cache-hierarchy/>

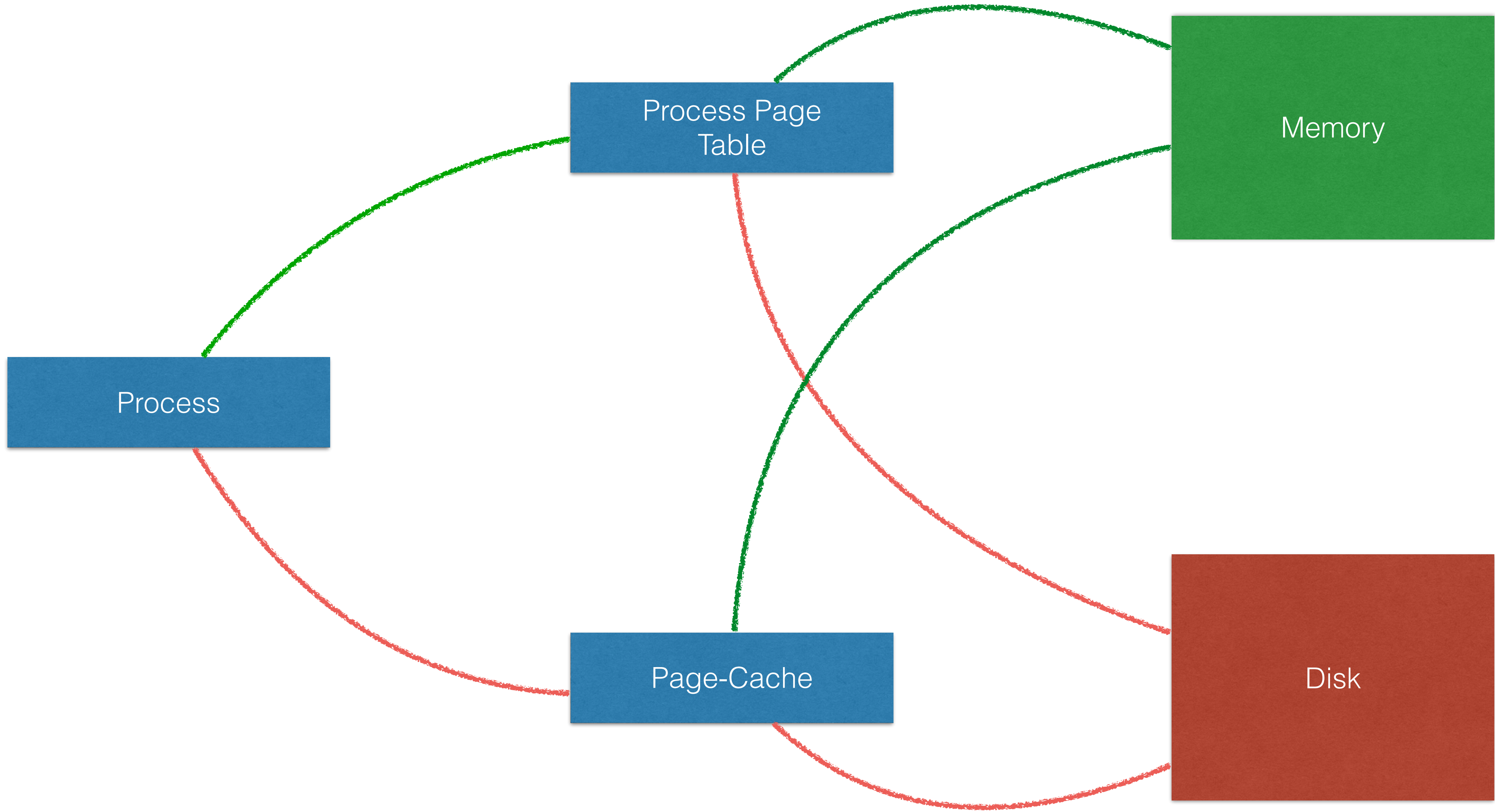
Is it faster to read from
a HashMap or a file?

Resident vs. Virtual Memory

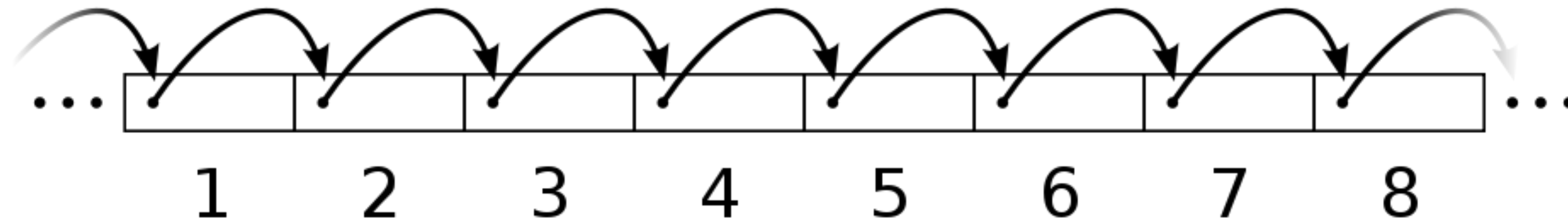
- How much memory can a process use?
 - OS maps when actually used
- What happens when we exceed physical?
 - Swap
 - Page faults



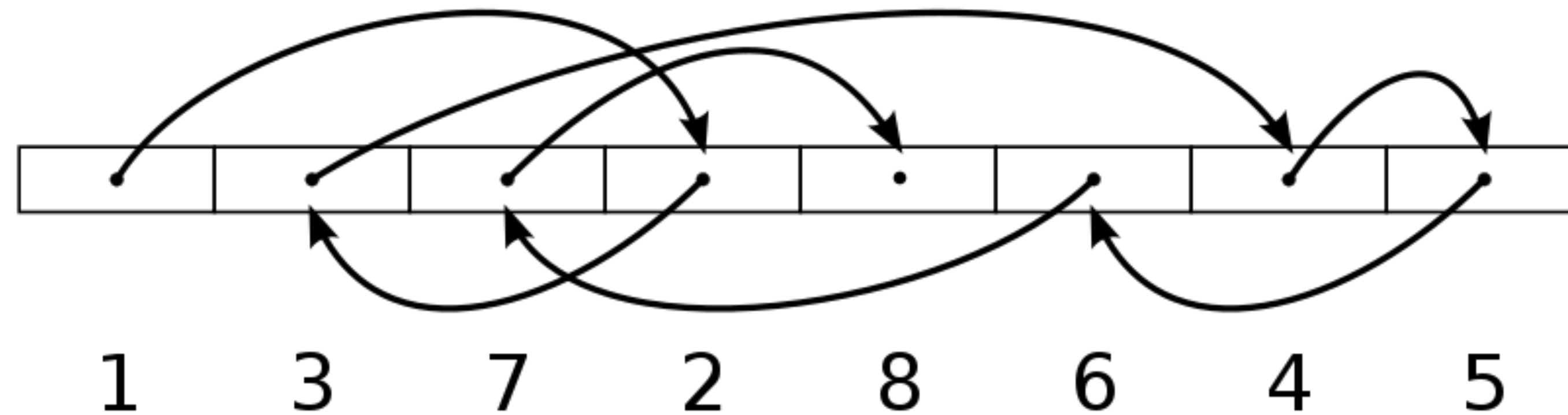
"Resident evil ver4" by impawards. Licensed under Fair use of copyrighted material in the context of Resident Evil (film) via Wikipedia - http://en.wikipedia.org/wiki/File:Resident_evil_ver4.jpg#mediaviewer/File:Resident_evil_ver4.jpg



Sequential access

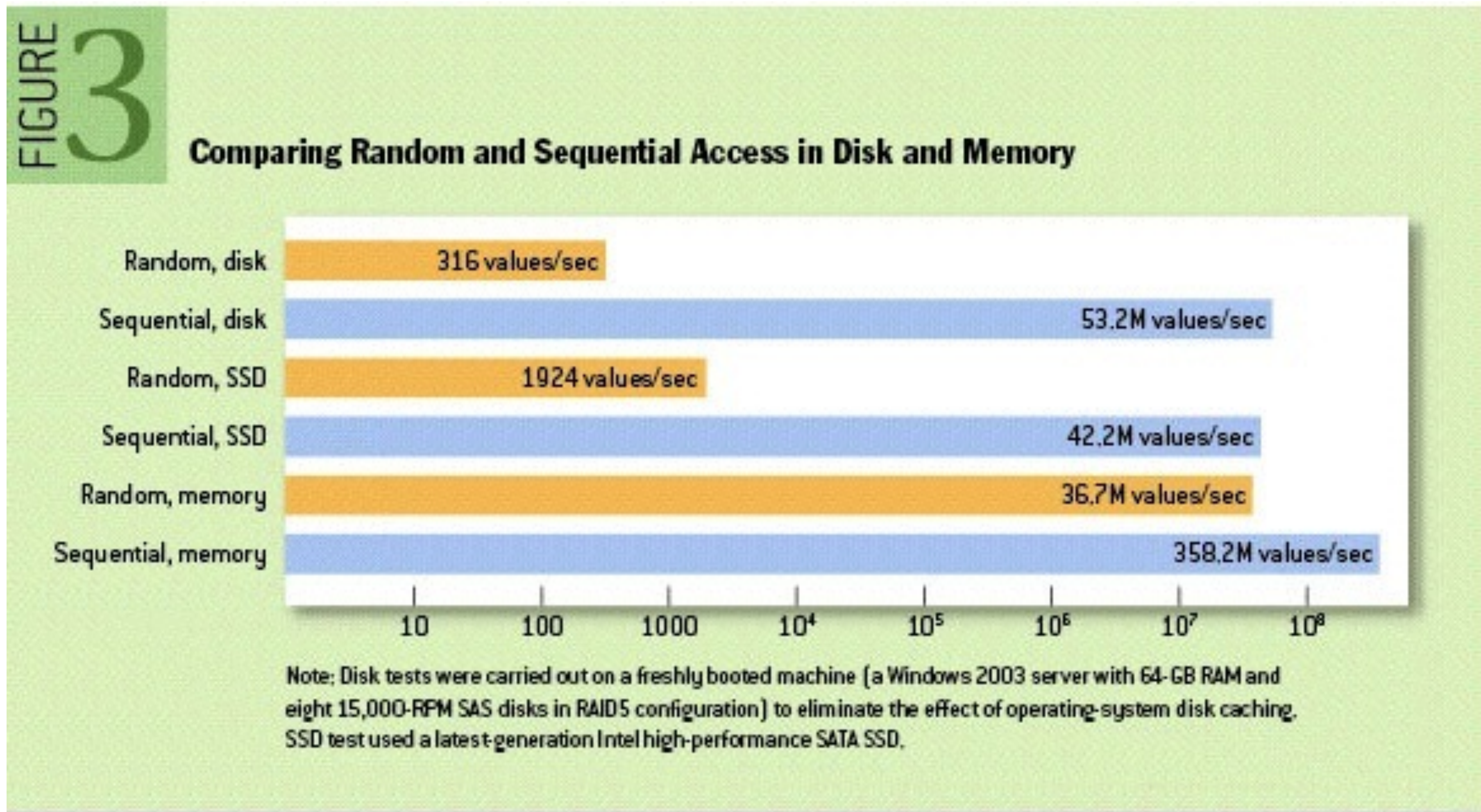


Random access



"Random vs sequential access" by (User:Intgr) - Own work. Licensed under Public Domain via Wikimedia Commons - http://commons.wikimedia.org/wiki/File:Random_vs_sequential_access.svg#mediaviewer/File:Random_vs_sequential_access.svg

Access Pattern Matters!



The Pathologies of Big Data - <http://queue.acm.org/detail.cfm?id=1563874>

So What?

- Benchmarking? Use relevantly sized data sets and access patterns
- Avoid random access if you can
- Avoid using more memory than available
- Performance counters(perf/likwid...):
 - Cache misses (L1/L2/LLC)
 - Page faults
- Disable swap? (Set `swappiness=0`)
- Consider priming memory?



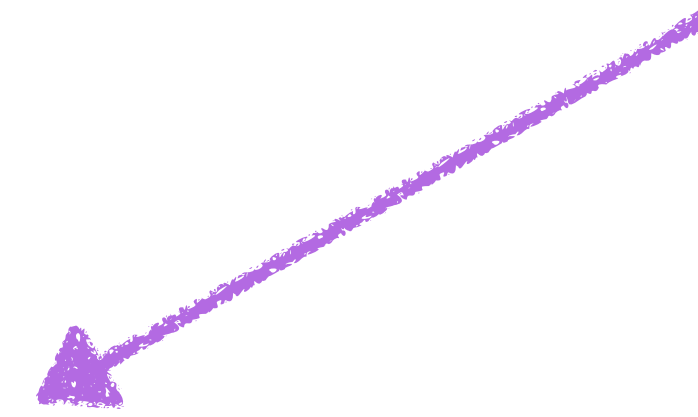
OOPs

hold = love;

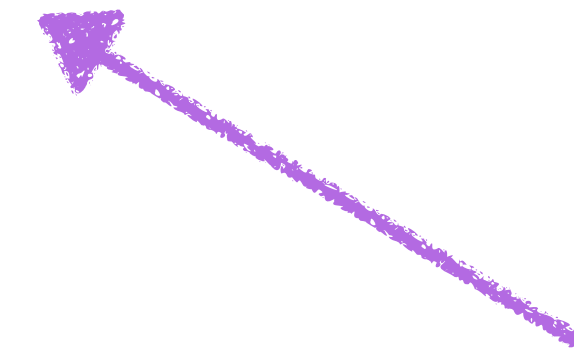


```
mov    r10,rsi
mov    r11,rdx
shr    r11,0x3
mov    DWORD PTR [rsi+0x10],r11d
shr    r10,0x9
movabs r11,0x10a4e9000
mov    BYTE PTR [r11+r10*1],r12b
```

Compressed Oops



Card Marking



What's an OOP?

- Ordinary Object Pointer
- Java: Object reference -> JVM: OOP
- Pointers to managed data on the heap

Memory Barrier

(not the JMM kind)

“...a block [of code, executed] on reading from or writing to certain memory locations by certain threads or processes.”

Memory Management Reference:

<http://www.memorymanagement.org/glossary/b.html#term-barrier-1>

Java Hidden Symbols

```
void copyPoint(Point p1, Point p2) {  
    p1.x = p2.x;  
}
```

```
void copyPoint(oop p1, oop p2) {  
    address a1 = readBarrier(p1);  
    address a2 = readBarrier(p2);  
    oop x = getObject(a2+xFieldOffset);  
    putObject(a1+xFieldOffset, x);  
    writeBarrier(a1, x);  
    safepoint_poll();  
}
```

Compressed OOPs

-XX:+UseCompressedOops

- Want large heaps (> 4Gb) and 32bit addresses
- Objects aligned to 8 bytes*
- Can compress OOP by dropping last 3 bits (>>3)
- Must decompress address to use it (<<3)
- Max referenced heap size is now 4Gb * 8 = 32Gb

<https://wikis.oracle.com/display/HotSpotInternals/CompressedOops>

* Can change with -XX:ObjectAlignmentInBytes

Compressed Oops Example(x86):

JAVA:

```
long v = this.l.longValue();
```

-XX:-UseCompressedOops:

```
mov r10, QWORD PTR [rsi+0x18]           ; r10= this.l  
mov r10, QWORD PTR [r10+0x10]         ; r10= l.value
```

-XX:+UseCompressedOops:

```
mov r11d, DWORD PTR [rsi+0x10]         ; r11d= this.l  
mov r10, QWORD PTR [r12+r11*8+0x10]; r10= l.value
```

CompressedOops -> Read Barrier

- Must be decompressed before read 'through'
- Can be copied without decompression
- Can be compared without decompression

LVB - Zing Read Barrier

- Looks like this

```
test8 rax, [gc_phase_trap_mask]; GC phase changed?
```

```
jnz 0x500d639b; GOTO LVB cold path
```

- Cold path: value has relocated
 - Mutator will fix up the loaded value
 - ‘Self healing’ - mutator participates in relocation
- Read more:

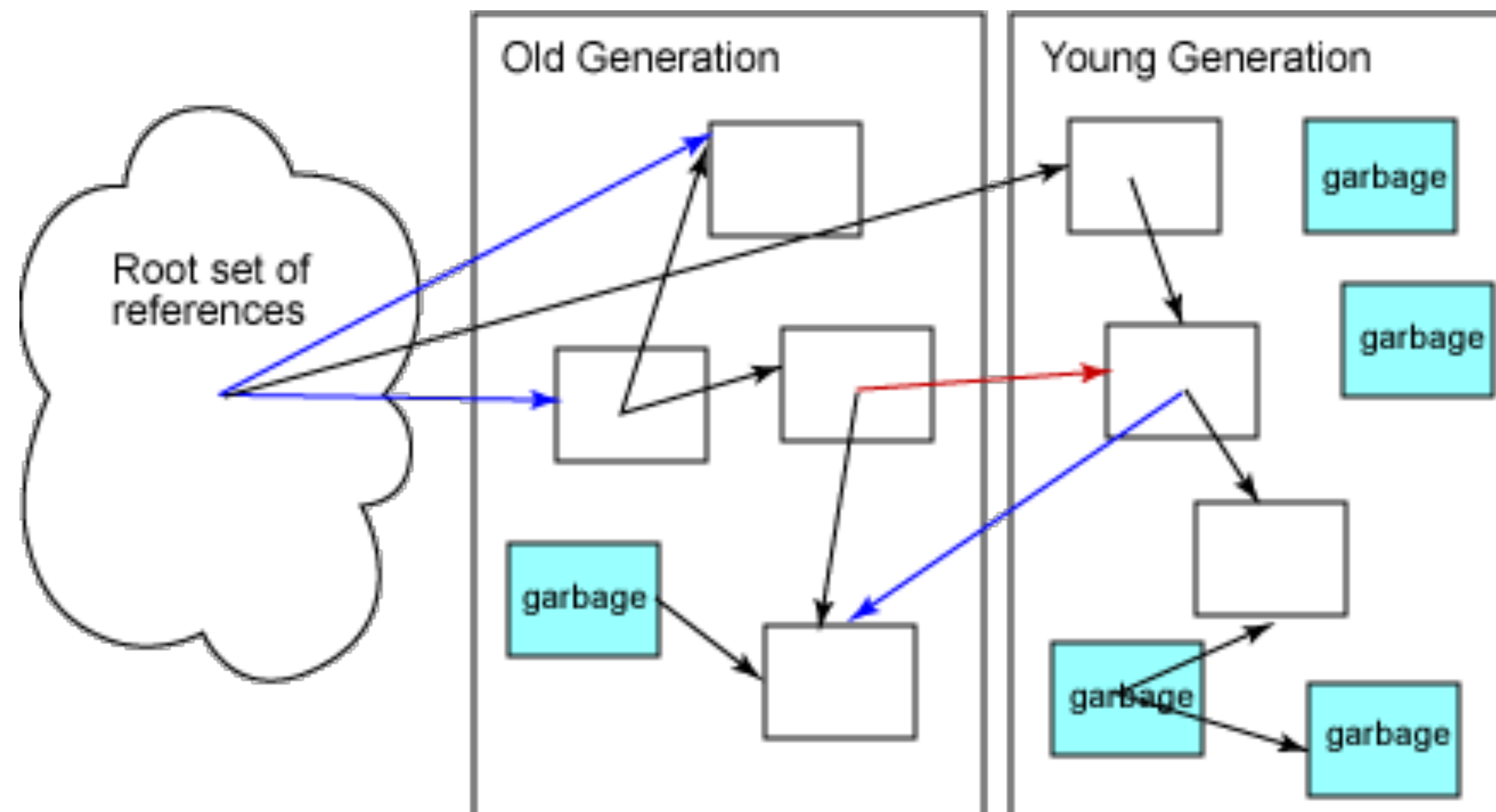
http://www.azulsystems.com/sites/default/files/images/c4_paper_acm.pdf

<http://www.javaworld.com/article/2078661>

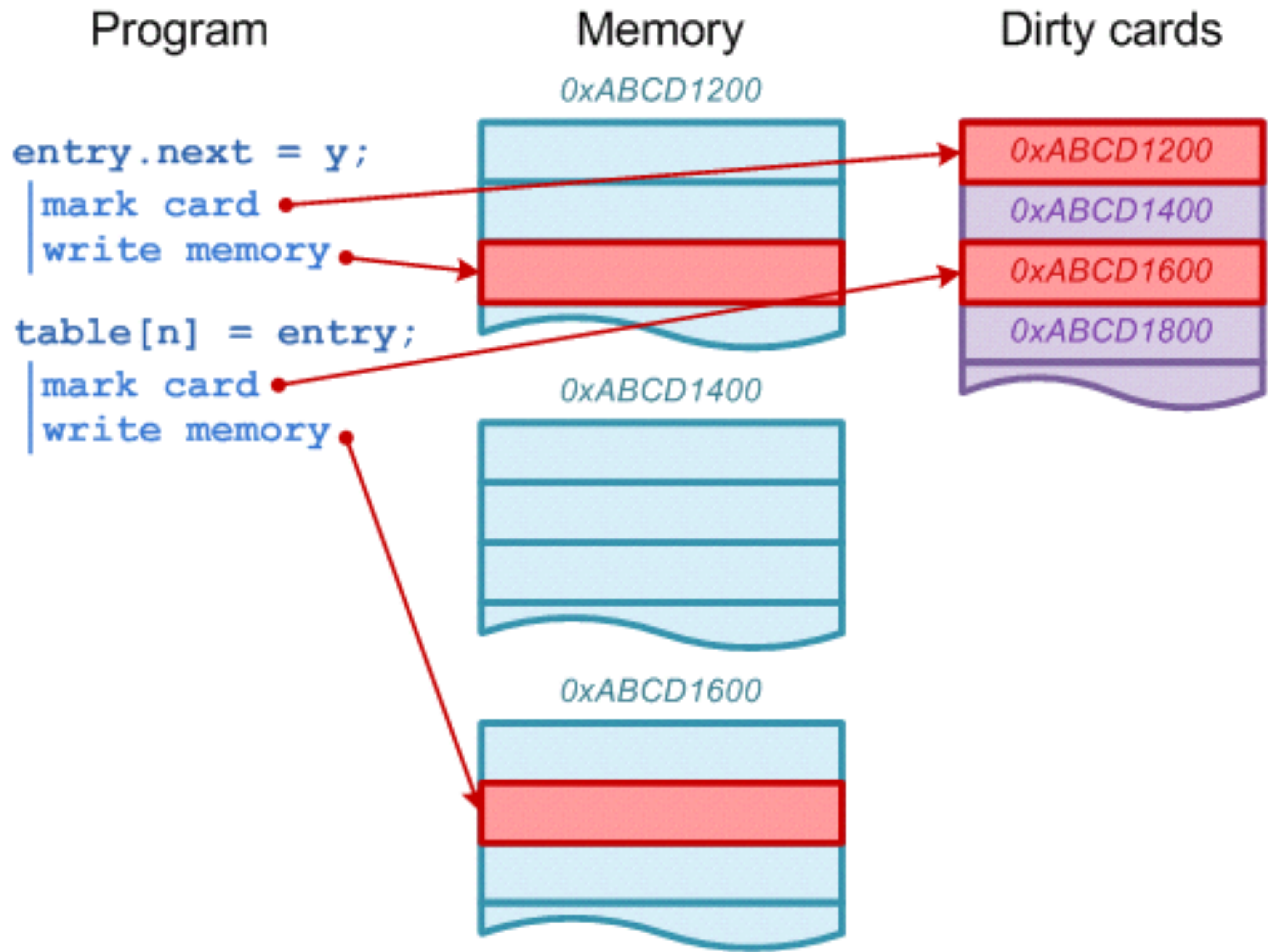


Card Marking

“The JVM maintains a card map, with one byte corresponding to each card in the heap. Each time a pointer field in an object in the heap is modified, the corresponding byte in the card map for that card is set.”



Brian Goetz - <http://www.ibm.com/developerworks/library/j-jtp11253/>



Card Marking is a Write Barrier

- An optimisation for young collections
- Reduce the impact of OldGen size on scan time
- Introduce a small overhead
- Comes in different flavours!

CardMarking v1 (default)

```
; rsi is 'this' address  
; rdx is setter param, reference to bar  
; this.foo = bar  
mov    QWORD PTR [rsi+0x20],rdx ; <- WHAT WE WANTED  
mov    r10,rsi  
; r10 = r10 >> 9;  
shr    r10,0x9  
; r11 is base of card table, byte[] CARD_TABLE  
mov    r11,0x7ebdfcff7f00  
; Mark 'this' card as dirty  
; CARD_TABLE[this address >> 9] = 0  
mov    BYTE PTR [r11+r10*1],0x0
```

CardMarking v2

(-XX:+UseCondCardMark)

```
; rsi is 'this' address
; rdx is setter param, reference to bar
; r10 = this
mov    r10,rsi    <- WHAT WE WANTED
; r10 = r10 >> 9
shr    r10,0x9
; r11 = CARD_TABLE
mov    r11,0x7f7cb98f7000
; r11 = CARD_TABLE + (this >> 9)
add    r11,r10
; r8d = CARD_TABLE[this >> 9]
movsx  r8d,BYTE PTR [r11]
test   r8d,r8d
; if(CARD_TABLE[this >> 9] == 0) goto 0x00007fc4a1071d7d
je     0x00007fc4a1071d7d
; CARD_TABLE[this >> 9] = 0
mov    BYTE PTR [r11],0x0
0x00007fc4a1071d7d:
mov    QWORD PTR [rsi+0x20],rdx ; this.foo = bar
```

CardMarking v3 (-XX:+UseG1GC)

```
movsx edi, BYTE PTR [r15+0x2d0] ; read GC flag
cmp edi, 0x0; if (flag != 0)
jne 0x00000001066fc601; GOTO OldValBarrier
Label WRITE:

mov QWORD PTR [rsi+0x20], rdx; this.foo = bar
mov rdi, rsi; rdi = this
xor rdi, rdx; rdi = this XOR bar
shr rdi, 0x14; rdi = (this XOR bar) >> 20
cmp rdi, 0x0; If this and bar are not same gen
jne 0x00000001066fc616; GOTO NewValBarrier
Label EXIT:
;...
Label OldValBarrier:
mov rdi, QWORD PTR [rsi+0x20]
cmp rdi, 0x0; if(this.foo == null)
je 0x00000001066fc5dd; GOTO WRITE
mov QWORD PTR [rsp], rdi ; setup rdi as parameter
call 0x000000010664bca0 ; {runtime_call}
jmp 0x00000001066fc5dd; GOTO WRITE
Label NewValBarrier:
cmp rdx, 0x0; bar == null
je 0x00000001066fc5f5 goto Exit
mov QWORD PTR [rsp], rsi
call 0x000000010664bda0 ; {runtime_call}
jmp 0x00000001066fc5f5 ; GOTO exit;
```

<- WHAT WE WANTED

So What?

- References mean extra work
- Impact can change by option/GC/JVM
- ‘Normalized/Flat’ data structures can help
 - Inheritance vs. Composition?
 - Consider access patterns
- Value Types might help (Java 9?)
- ObjectLayout might help (Java 9?)



The JVM...



... Must give
us pause?

Why Stop The World?

- Some GC phases
- Deoptimization
- Stack trace dump
- Lock un-biasing
- Class redefinition

Can one Java thread stop
all other Java threads?

You too can trigger a STW pause!

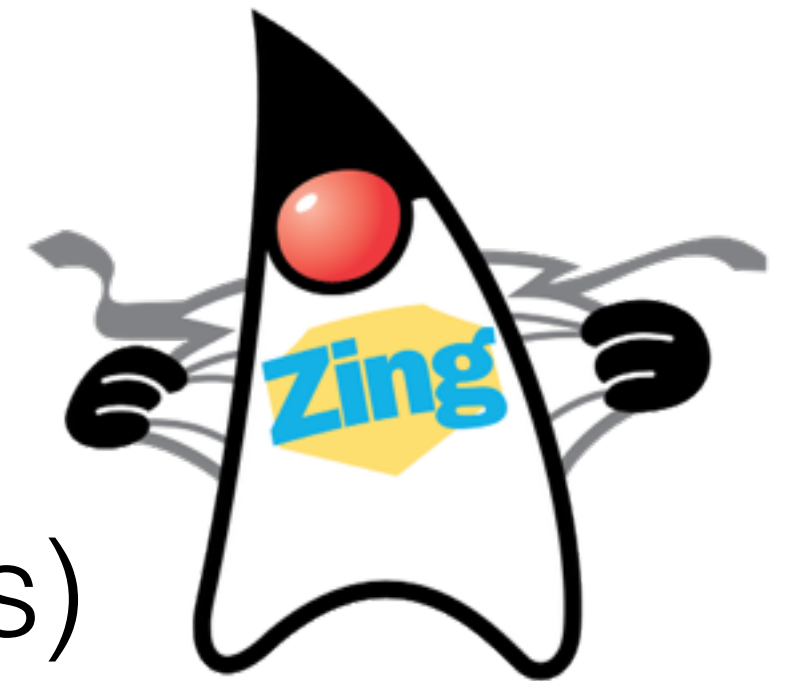
- On normal allocation (Young Gen exhausted*)
- On large object allocation (Old Gen exhausted*)
- On synchronized block (unbiasing*)
- Profiler sampling (JVMTI *GetStackTrace()**)
- Hitting cold code (deoptimization**)

* Not true for Zing

** Less true for Zing

Zing: The pause-**less** JVM

- C4 - Fully Concurrent GC (including young gen)
- ReadyNow! - Fighting Deopts
- No biased locking
- ... some STW pauses exist (smaller than OS hiccups)



To 'Stop The World' the
JVM brings all threads to a
SAFEPPOINT

- Raise Safepoint 'flag'
- Wait for ALL threads to reach Safepoint state

Safepoint

(noun.)

- A thread state
 - `Waiting/Idle/Blocked` -> `@Safepoint`
 - `Running/Ready Java` -> `!@Safepoint`
 - `Running/Ready native` -> `@Safepoint`

<http://blog.ragozin.info/2012/10/safepoints-in-hotspot-jvm.html>

<http://psy-lob-saw.blogspot.com/2014/03/where-is-my-safepoint.html>

<http://chriskirk.blogspot.ru/2013/09/what-is-java-safepoint.html>

Safepoint poll: OpenJDK

- Read from a special page:

```
test    DWORD PTR [rip+0xffffffffffe690e53],eax
```

- JVM Sets the page to protected, polling threads trap a SEGV and go to safepoint
- Look for `{poll}` or `{poll_return}` in the assembly comments

Safepoint poll: Zing

- Read the thread local safe point flag:

```
gs:cmp4i [0x40 t1s._please_self_suspend],0
```

```
jnz 0x500a0186; Where the safepoint code be
```

- JVM Sets the thread flag to 1, polling threads hop to
- Look for `t1s._please_self_suspend`

Where do we see a Safepoint poll?

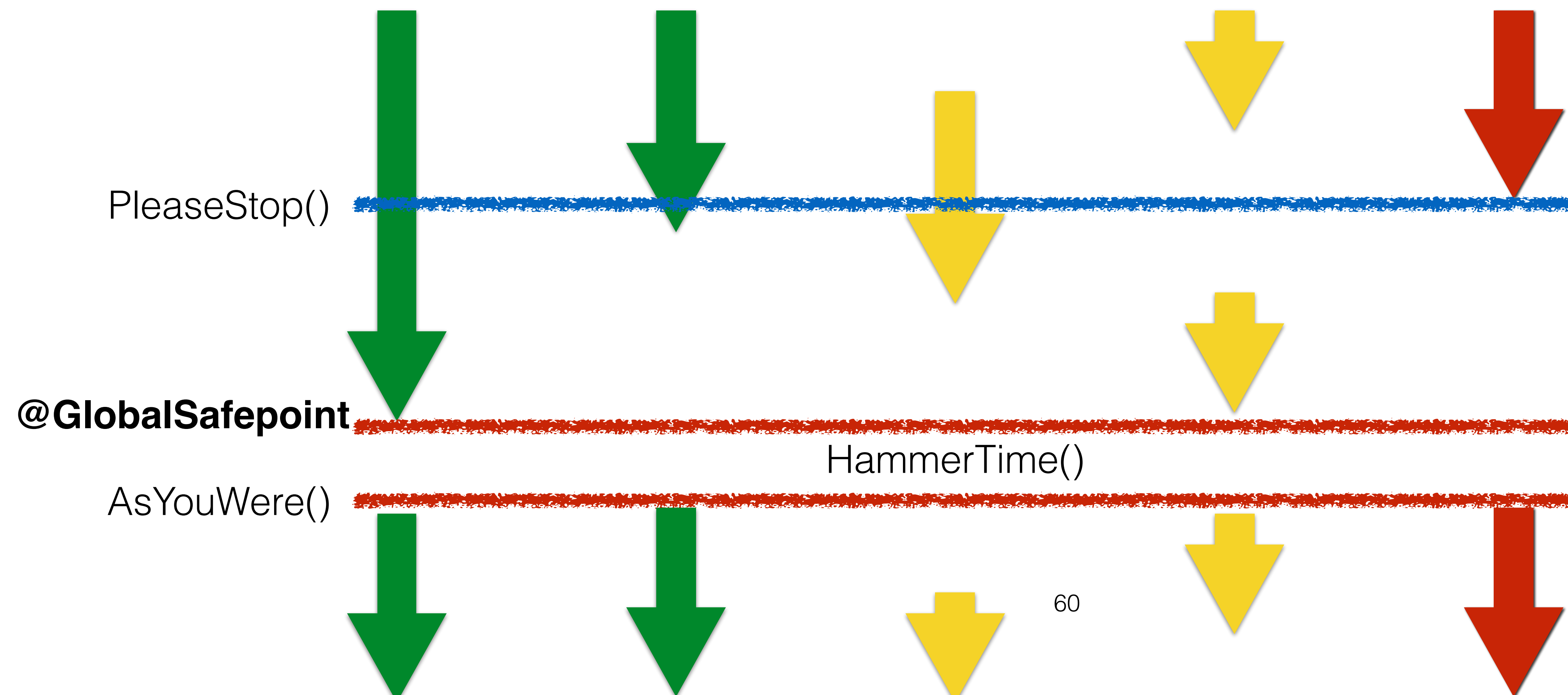
- While loop back edge
- For loop over longs back edge (for(long l=0;...))
- Method exit

How long was this GC pause?

```
8109.128: [GC [PSYoungGen: 109884K-  
>14201K(139904K) ] 691015K->595332K(1119040K) ,  
0.0454530 secs]
```

TTSP - Time To Safepoint

- Not included in GC Time (use `-XX:+PrintGCApplicationStoppedTime`)
- Max Actual Pause = TTSP + pause time(in safepoint)



What can cause long TTSP?

- Inlining removes end of method safe point poll
- Long counted loops* (i.e: `for(int i=0;i<MAX_INT;i++) doSomething();`)
- Large memory copies (`System.arraycopy`**/`Unsafe.copyMemory`)
- Interrupted threads
- Page Faults

* On Zing can use: `-XX+KeepSafepointsInCountedLoops`

** On Zing safepoint polls are introduced between copying chunks

So what?

- Use `-XX:+PrintGCApplicationStoppedTime`
- TTSP can dominate stop the world times
- Beware Safepoint biased profilers
- Chunk large memory copies



```
while (hasQ() && hasTime()) {  
    A();  
}  
return;
```